Improving Parallel Scalability of Spectral Element Method Pressure Poisson Preconditioners (With 4th-kind Chebyshev Smoothing)

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# **Big Picture**



Figure: 352K pebble geometry from<sup>1</sup>, n = 51B, P = 27648 V100s on Summit.

<sup>&</sup>lt;sup>1</sup>Min, Lan, Fischer, Merzari, Kerkemeier, Phillips, Rathnayake, Novak, Gaston, Chalmers, et al., "Optimization of full-core reactor simulations on summit", 2022.

### Motivation



nekRS Timing Breakdown: n=51B, 2000 Steps					
	pre-tun	ing	post-tuning		
Operation	time (s)	%	time (s)	%	
computation	1.19 + 03	100	5.47+02	100	
advection	5.82 + 01	5	4.49 + 01	8	
viscous update	5.38+01	5	5.98+01	. 11	
pressure solve	1.08 + 03	90	4.39+02	80	
precond.	9.29+02	78	3.67+02	67	
coarse grid	5.40+02	45	6.04 + 01	11	
projection	6.78+00	1	1.21 + 01	2	
dotp	4.92+01	4	1.92 + 01	4	

Table: Runtime statistics for the 352K pebble geometry of fig. 1 on P = 27648 V100s on Summit.

#### Poisson

Solve series of Poisson problems using SE discretization:

$$-\nabla^2 \tilde{u} = \tilde{f} \text{ for } \tilde{u}, \tilde{f} \in \Omega \subset \mathbb{R}^d \mapsto \mathbb{R}.$$
(1)

Weak formulation: find  $u^m(\mathbf{x}) \in X_0^N \subset \mathcal{H}_0^1$  such that

$$\int_{\Omega} \nabla v \cdot \nabla u^m \, dV = \int_{\Omega} v \, f^m \, dV \quad \forall \, v \in X_0^N, \qquad (2)$$

$$X_0^N = \operatorname{span}\{\phi_j(\mathbf{x})\}$$
(3)

Discrete problem – solve  $A\underline{u}^m = \underline{b}^m$ :

$$a_{ij} = \int_{\Omega} \nabla \phi_i \cdot \nabla \phi_j \, dV. \tag{4}$$

How to solve? Multigrid.

#### **Polynomial Smoothers**

Polynomial smoother  $G_j = (I - \omega S_j A_j)^k$  is k steps of simple smoothing iteration:

$$\left(\underline{x}_{i+1}\right)_j = (\underline{x}_i)_j + \omega S_j(\underline{b}_j - A_j(\underline{x}_i)_j).$$
(5)

Can we do better?

$$\min_{p_k \in \mathbb{P}_k, p_k(0)=1} \max_{\lambda \in [\lambda_{\min}, \lambda_{\max}]} |p(t)|.$$
(6)

#### 1st-kind Chebyshev Smoother<sup>23</sup>

Minimax solution:

$$\hat{T}_k(\lambda) = \frac{1}{\sigma_k} T_k\left(\frac{\theta - \lambda}{\delta}\right) \text{ with } \sigma_k := T_k\left(\frac{\theta}{\delta}\right).$$
 (7)

 $T_k$  are Chebyshev polynomials of the 1st-kind:

$$T_n(x) = 2xT_{n-1}(x) - T_{n-2}(x)$$
  
 $T_0(x) = 1$   
 $T_1(x) = x.$  (8)

 $\theta$  is the midpoint of the interval  $[\lambda_{\min}, \lambda_{\max}]$ :

$$\theta = \frac{\lambda_{\min} + \lambda_{\max}}{2}.$$

 $\delta$  is the mid-width of the interval:

$$\delta = rac{\lambda_{max} - \lambda_{min}}{2}$$

<sup>2</sup>Adams, Brezina, Hu, and Tuminaro, "Parallel multigrid smoothing: polynomial versus Gauss-Seidel", 2003.

<sup>3</sup>Kronbichler and Ljungkvist, "Multigrid for matrix-free high-order finite element computations on graphics processors", 2019.



Figure: Smoother polynomials for the simple smoother (a) and the 1st-kind Chebyshev smoother (b).

Algorithm Chebyshev smoother, 1st-kind

$$\theta = \frac{1}{2} (\lambda_{max} + \lambda_{min}), \ \delta = \frac{1}{2} (\lambda_{max} - \lambda_{min}), \ \sigma = \frac{\theta}{\delta}, \ \rho_0 = \frac{1}{\sigma}$$

$$\underline{x}_0 = \underline{x}, \underline{r}_0 = S(\underline{b} - A\underline{x}_0), \ \underline{d}_0 = \frac{1}{\theta}\underline{r}_0$$
for  $i = 1, \dots, k - 1$  do
$$\underline{x}_i = \underline{x}_{i-1} + \underline{d}_{i-1}$$

$$\underline{r}_i = \underline{r}_{i-1} - SA\underline{d}_{i-1}, \ \rho_i = \frac{1}{2\sigma - \rho_{i-1}}$$

$$\underline{d}_i = \rho_i\rho_{i-1}\underline{d}_{i-1} + \frac{2\rho_i}{\delta}\underline{r}_i$$
end for
$$\underline{x}_k = \underline{x}_{k-1} + \underline{d}_{k-1}$$
return  $\underline{x}_k$ 

### 4th-kind Chebyshev Smoother<sup>5</sup>

wlog  $\rho(SA) = 1$ . *Two-level* Hackbusch bound<sup>4</sup>:

$$\|E_{\searrow}\|_{A} = \left\| (I - PA_{c}^{-1}P^{T}A)G_{k} \right\|_{A}$$
$$\leq C^{1/2} \sup_{0 < \lambda \leq 1} \lambda^{1/2} |p_{k}(\lambda)|.$$
(9)

What  $p_k$  minimizes this error bound?

<sup>&</sup>lt;sup>4</sup>Hackbusch, "Multi-grid convergence theory", 1982.

<sup>&</sup>lt;sup>5</sup>Lottes, "Optimal polynomial smoothers for multigrid V-cycles", 2022.

Weighted minimax solution:

$$p_k(\lambda) = \frac{1}{2k+1} W_k(1-2\lambda),$$
 (10)

 $W_k$  are 4th-kind Chebyshev polynomial<sup>6</sup>:

$$W_n(x) = 2xW_{n-1}(x) - W_{n-2}(x)$$
  
 $W_0(x) = 1$   
 $W_1(x) = 2x + 1.$  (11)

Can we do even better? What about the multi-level case?

<sup>&</sup>lt;sup>6</sup>Mason, "Chebyshev polynomials of the second, third and fourth kinds in approximation, indefinite integration, and integral transforms", 1993.

Lemma<sup>7</sup>:

Let the smoother iteration (on each level j) be given by

$$G_j = p_{k_j}(S_j A_j)$$

where  $S_j$  is SPD,  $\rho(S_jA_j) = 1$ , and  $p_{k_j}(x)$  is a  $k_j$ -order polynomial satisfying  $p_{k_j}(0) = 1$  and  $|p_{k_j}(x)| < 1$  for  $0 < x \le 1$ , possibly different on each level. Then the V-cycle contraction factor

$$\|E_{\nearrow}\|_{A}^{2} \leq \max_{j \in 0, \dots, \ell-1} \frac{C_{j}}{C_{j} + \gamma_{j}^{-1}}$$
 (12)

where  $C_j$  is the approximation property constant for level j, and

$$\gamma_j = \sup_{0 < \lambda \le 1} \frac{\lambda \, p_{k_j}(\lambda)^2}{1 - p_{k_j}(\lambda)^2}.$$
(13)

<sup>&</sup>lt;sup>7</sup>Lottes, "Optimal polynomial smoothers for multigrid V-cycles", 2022.



Figure: 4th-kind Chebyshev smoother (a) and the 4th-kind Chebyshev smoother optimized with respect to previous error bound (b).



Figure: All smoother polynomials at order 3.

$$\|E_{\nearrow}\|_A^2 \leq \frac{C}{C+\gamma^{-1}}$$

Polynomial Smoother	$ig  \gamma^{-1}, k  o \infty$
Simple multi-sweep, damping	$ $ $2\omega k$
1st-kind Chebyshev, fixed $\lambda_{min}$	$2\sqrt{\frac{1}{\lambda_{min}}}k$
1st-kind Chebyshev, $\lambda^*_{\min}$ optimizes $\gamma^{-1}$	$ $ 2.38 $k^{1.78}$
4th-kind Chebyshev	$\left  \frac{4}{3}k(k+1)\right $
4th-kind optimal Chebyshev	$\left  \begin{array}{c} \frac{4}{\pi^2}(2k+1)^2 - \frac{2}{3} \end{array} \right $

$$\begin{split} \underline{x}_0 &= \underline{x}, \ \underline{r}_0 = \underline{b} - A\underline{x}_0\\ \underline{d}_0 &= \frac{4}{3} \frac{1}{\lambda_{max}} S\underline{r}_0\\ \text{for } i &= 1, \dots, k-1 \text{ do}\\ \underline{x}_i &= \underline{x}_{i-1} + \beta_i \underline{d}_{i-1}, \ \underline{r}_i &= \underline{r}_{i-1} - A\underline{d}_{i-1}\\ \underline{d}_i &= \frac{2i-1}{2i+3} \underline{d}_{i-1} + \frac{8i+4}{2i+3} \frac{1}{\lambda_{max}} S\underline{r}_i\\ \text{end for}\\ \underline{x}_k &= \underline{x}_{k-1} + \beta_k \underline{d}_{k-1}\\ \text{return } \underline{x}_k \end{split}$$

<u>Used for *p*-multigrid (pMG)</u> and algebraic multigrid (AMG) <sup>8</sup>. <sup>8</sup>AMG implementations available for:

- hypre/boomerAMG: https://github.com/MalachiTimothyPhillips/ hypre/tree/fourth-kind-chebyshev-polynomials
- Trilinos/MueLu: https://github.com/MalachiTimothyPhillips/ Trilinos/tree/optimal-chebyshev-polynomials

#### To Post-smooth, or Not to Post-smooth?

- Given 2k smoother passes, what order m pre-smoothing, n post-smoothing should be used, m + n = 2k?
- Answer using error bound from previous Lemma:

$$\arg \max_{m,n,m+n=2k} C \left( \gamma^{-1}(m) + \gamma^{-1}(n) \right) + \gamma^{-1}(m) \cdot \gamma^{-1}(n) \quad (14)$$

- Check solutions for k < 50 in SymPy<sup>9</sup>.
- With few exceptions, either m = n = k (symmetric smoothing) or m = 2k, n = 0 (no post-smoothing) is optimal<sup>10</sup>.
- When to use which?

<sup>9</sup>Meurer, Smith, Paprocki, Čertík, Kirpichev, Rocklin, Kumar, Ivanov, Moore, Singh, et al., "SymPy: symbolic computing in Python", 2017. <sup>10</sup>Phillips and Fischer, "Optimal Chebyshev Smoothers and One-sided V-cycles", 2022.

### To Post-smooth, or Not to Post-smooth?

Polynomial Smoother	When to <i>omit</i> post smoothing?
Simple multi-sweep, damping	$C > \frac{(4k - \log(4k))^2}{\log(2k)}$
1st-kind Chebyshev, fixed $\lambda_{min}=0.1$	$\mid$ $k>$ 3, $C \gtrsim 1.55 e^{1.45k}$
1st-kind Chebyshev, $\lambda_{\min}^*$ optimizes $\gamma^{-1}$	$C \gtrsim 2.38 k^{1.78}$
4th-kind Chebyshev	$C > \frac{2(k+1)^2}{3}$
4th-kind optimal Chebyshev	$  C > \frac{2(6(2k+1)^2 - \pi^2)^2}{3\pi^2(-12(2k+1)^2 + 6(4k+1)^2 + \pi^2)}$

C is the multigrid approximation property constant. Roughly  $\kappa(SA)$  restricted to the A-orthogonal complement of the coarse-grid space.

## nekRS Pressure Poisson Results

Solver parameter study in nekRS:

- Consider 3 smoothers for Chebyshev-acceleration:
  - Jacobi
  - Additive Schwarz Method (ASM)
  - Restrictive Additive Schwarz (RAS)
- Consider 4 types of polynomial acceleration schemes:
  - 1st-kind Chebyshev
  - 1st-kind Chebyshev,  $\lambda_{\min}$  optimized via random RHS
  - 4th-kind Chebyshev<sup>11</sup>
  - Optimized 4th-kind Chebyshev
- Vary k from 1 to 6
- Consider 2 different V-cycle approaches:
  - (k, k) symmetric V-cycle
  - (2k,0) V-cycle (no post-smoothing)

<sup>&</sup>lt;sup>11</sup>Lottes, "Optimal polynomial smoothers for multigrid V-cycles", 2022.

## nekRS Pressure Poisson Results



Case Name	E	N	n
146 pebble (fig. 8a)	62K	7	21M
1568 pebble (fig. 8b)	524K	7	180M
67 pebble (fig. 8c)	122K	7	42M
Kershaw ( $\varepsilon = 1.0$ ) (fig. 8d)	47K	7	16M
Kershaw ( $\varepsilon = 0.3$ ) (fig. 8e)	47K	7	16M
Kershaw ( $\varepsilon = 0.05$ ) (fig. 8f)	47K	7	16M

Table: Discretization and fastest solver for the NS examples.



Case	Ρ	Fastest Solver	Ts	lter.	$\frac{T_D}{T_S}$	$\frac{(T_{crs})_D}{(T_{crs})_S}$
(a) pb146	6	4 <sup>th</sup> <sub>opt</sub> -Cheb, RAS(4,4)	0.15	5.3	1.17	1.21
(b) pb67	18	$4_{opt}^{th}$ -Cheb, RAS(12,0)	0.37	12.5	1.81	2.41
(c) pb1568	72	4 <sup>th</sup> -Cheb, ASM(12,0)	0.14	3	1.27	2.13
(d) K. 1	6	1 <sup>st</sup> -Cheb, $\lambda_{min}^{opt}$ , RAS(2,2)	0.09	8	1.75	1.13
(e) K. 0.3	6	$1^{st}$ -Cheb, $\lambda_{min}^{opt}$ , RAS(5,5)	0.67	28	1.35	1.79
(f) K. 0.05	6	$4_{opt}^{th}$ -Cheb, RAS(12,0)	2.40	88	1.75	2.31

Table:  $T_S$ : solution time of fastest solver.  $T_D$  solution time of nekRS default, 1<sup>st</sup>-Cheb, ASM(3,3).

## Questions?

- More details in pre-print: "Optimal Chebyshev Smoothers and One-sided V-cycles" https://arxiv.org/abs/2210.03179
- nekRS: https://github.com/Nek5000/nekRS
- 4th-kind Chebyshev implementations in popular AMG solvers:
  - hypre/boomerAMG: https://github.com/MalachiTimothyPhillips/ hypre/tree/fourth-kind-chebyshev-polynomials
  - Trilinos/MueLu: https://github.com/MalachiTimothyPhillips/ Trilinos/tree/optimal-chebyshev-polynomials

Supporting Materials

## **Operator Cost**



Figure: Weak scaling operator cost study for the Poisson solver for the Kershaw benchmark problem, n/P = 2.67M,  $\varepsilon = 0.05^{-12}$ .

<sup>12</sup>log(*P*) scaling of coarse grid solve, SEMFEM operator are expected, see Fischer, "Scaling limits for PDE-based simulation", 2015; Tufo and Fischer, "Fast parallel direct solvers for coarse grid problems", 2001.

Opt. 4th-kind Chebyshev polynomial for  $\underline{e}_k = p_k(SA)\underline{e}_0$ :

$$p_k(\lambda) = \sum_{i=0}^k \frac{\beta_i - \beta_{i+1}}{2i+1} W_i(1-2\lambda),$$
 (15)

with  $\beta_0 = 1$  and  $\beta_{k+1} = 0$ .

## Multigrid approximation property constant

$$C_{j} := \left\| A_{j}^{-1} - P_{j+1}^{j} A_{j+1}^{-1} \left( P_{j+1}^{j} \right)^{T} \right\|_{A_{j},S_{j}}^{2}$$
$$:= \sup_{\|\underline{f}\|_{S_{j}} \leq 1} \left\| \left( A_{j}^{-1} - P_{j+1}^{j} A_{j+1}^{-1} \left( P_{j+1}^{j} \right)^{T} \right) \underline{f} \right\|_{A_{j}}^{2}.$$
(16)

 $C_j$  is roughly  $\kappa(S_jA_j)$  restricted to the  $A_j$ -orthogonal complement of the coarse (j + 1)-space.

# *p*-multigrid

- Matrix-free a must:
  - dofs:  $n \sim Ep^3$
  - nnz(A)  $\sim O(Ep^6)$
  - $A\underline{x} \operatorname{cost} O(Ep^4) = O(np)$
- Drop Galerkin requirement for coarser levels
  - Each MG level has different polynomial order

• e.g., 
$$p = 7$$
,  $p = 3$ ,  $p = 1$ 

#### Schwarz-based Smoothers

SE-based additive Schwarz method (ASM) smoothers<sup>13</sup>:

$$S_{ASM\underline{r}} = \sum_{e=1}^{E} W_e R_e^T \bar{A}_e^{-1} R_e \underline{r}$$
(17)

Or, restrictive additive Schwarz  $(RAS)^{14}$ :

$$S_{RAS\underline{r}} = \sum_{e=1}^{E} \tilde{R}_{e}^{T} \bar{A}_{e}^{-1} R_{e\underline{r}}.$$
(18)

Subdomains area extensions of element with  $\bar{p}^3 = (p+3)^3$  dofs.

$$\bar{A}_e \neq R_e^T A_e R_e$$
 ruins  $O(pn)$  complexity (19)

<sup>13</sup>Lottes and Fischer, "Hybrid multigrid/Schwarz algorithms for the spectral element method", 2005; Loisel, Nabben, and Szyld, "On hybrid multigrid-Schwarz algorithms", 2008.

<sup>14</sup>Cai and Sarkis, "A restricted additive Schwarz preconditioner for general sparse linear systems", 1999.



Figure: Figure 6a Approximation of deformed elements  $\Omega_1$  and  $\Omega_2$  as box-shaped, overlapping subdomains  $\overline{\Omega}_1$  and  $\overline{\Omega}_2$ . Figure 6b overlapping subdomain  $\widetilde{\Omega}_e$ , constructed by overlapping two nodes in each spatial dimension and applying a homogeneous Dirichlet boundary condition on  $\partial \overline{\Omega}_e$ .

### Fast Diagonalization Method

$$ar{A}_e = B_z \otimes B_y \otimes A_x + B_z \otimes A_y \otimes B_x + A_z \otimes B_y \otimes B_x,$$

Generalized eigenvalue problem in x, y, z:

$$A_*\underline{s}_i = \lambda_i B_*\underline{s}_i$$

Fast, direct inverse:

$$\bar{A}_{e}^{-1} = (S_{z} \otimes S_{y} \otimes S_{x})D^{-1}(S_{z}^{T} \otimes S_{y}^{T} \otimes S_{x}^{T}),$$

$$D := I \otimes I \otimes \Lambda_x + I \otimes \Lambda_y \otimes I + \Lambda_z \otimes I \otimes I.$$

- Storage:  $3E\bar{p}^2 + E\bar{p}^3$
- Complexity:  $O(E\bar{p}^4)$
- Use Schwarz-based smoothers in Chebyshev acceleration

# Preconditioning via Low-order Operator

- Precondition high-order system using low-order system
- Spectral equivalence  $\kappa(A_F^{-1}A) \sim \pi^2/4$  in certain cases<sup>15</sup>
- Choice of finite element space matters<sup>16</sup>
  - Strong diagonal preconditioner,  $M^{-1} = A_F^{-1} B_d B^{-1}$ .
- Bello-Maldonado and Fischer<sup>17</sup> proposed one-per-vertex scheme
  - Use this with weak preconditioner,  $M^{-1} = A_F^{-1}$ .



<sup>15</sup>Orszag, "Spectral methods for problems in complex geometrics", 1979. <sup>16</sup>Canuto, Gervasio, and Quarteroni, "Finite-element preconditioning of G-NI spectral methods", 2010.

<sup>17</sup>Bello-Maldonado and Fischer, "Scalable low-order finite element preconditioners for high-order spectral element Poisson solvers", 2019.

- How to apply  $A_F^{-1}$ ? AMG!
  - PMIS coarsening
  - 0.25 strength threshold
  - Extended + i interpolation ( $p_{max} = 4$ )
  - L<sub>1</sub>-Jacobi relaxation
  - One V-cycle for preconditioning
  - Smoothing on the coarsest level
- Use either  $AmgX^{18}$  or boomerAMG<sup>19</sup> on GPU.
- Other approaches exist: Pazner, "Efficient low-order refined preconditioners for high-order matrix-free continuous and discontinuous Galerkin methods", 2020

<sup>&</sup>lt;sup>18</sup>Naumov, Arsaev, Castonguay, Cohen, Demouth, Eaton, Layton, Markovskiy, Reguly, Sakharnykh, et al., "AmgX: A library for GPU accelerated algebraic multigrid and preconditioned iterative methods", 2015.

<sup>&</sup>lt;sup>19</sup>Falgout, Li, Sjögreen, Wang, and Yang, "Porting hypre to heterogeneous computer architectures: Strategies and experiences", 2021.

Kershaw,  $\varepsilon = 0.05$ 



Figure: Weak scaling results for Kershaw,  $\varepsilon = 0.05$ .



Figure: Navier-Stokes cases: pebble-beds with (a) 146, (b) 1568, and (c) 67 spheres; (d) Boeing speed bump.

Case Name	E	Ν	n	Fastest Solver
146 pebble (fig. 8a)	62K	7	21M	1 <sup>st</sup> Cheb-RAS(3,3),(7,5,3,1)
1568 pebble (fig. 8b)	524K	7	180M	SEMFEM
67 pebble (fig. 8c)	122K	7	42M	SEMFEM (4X Speedup)
Speed bump (fig. 8d)	885K	9	645M	1 <sup>st</sup> Cheb-RAS(3,3),(9,5,1)

Table: Discretization and fastest solver for the NS examples.

#### 1st-kind Chebyshev

Correlation for  $\lambda^*_{\min}$  with 1% relative error and 0.1% absolute error for  $k \in [1, 50]$  is given by

$$\lambda_{\min}^* \approx \frac{1.69}{k^{1.68} + 2.11k + 1.98}.$$
 (20)